

An introduction to graph C^* -algebras

Lecture 5

David Pask

School of Mathematics and Applied Statistics
Wollongong University
AUSTRALIA

10th. July, 2019

Shift spaces

- Let \mathcal{A} be a set, usually finite, of symbols, called the *alphabet*.
- Let \mathcal{A}^* denote the Kleene Star of \mathcal{A} , which consists of *words* or finite strings of symbols drawn from \mathcal{A} .
- The *full \mathcal{A} -shift* is the set $\mathcal{A}^{\mathbb{Z}}$ together with the shift map $\sigma : \mathcal{A}^{\mathbb{Z}} \rightarrow \mathcal{A}^{\mathbb{Z}}$ defined by

$$(\sigma x)_n = x_{n+1} \quad n \in \mathbb{Z}.$$

- A *shift space* is a closed, shift invariant subset of $\mathcal{A}^{\mathbb{Z}}$.
- A shift space $X \subset \mathcal{A}^{\mathbb{Z}}$ can be described by giving a list $\mathcal{F} \subset \mathcal{A}^*$ of *forbidden words* that is

$$X = X_{\mathcal{F}} := \{y \in \mathcal{A}^{\mathbb{Z}} : \text{no subword of } y \text{ contains an element of } \mathcal{F}\}.$$

- The shift spaces (X_1, σ_1) , (X_2, σ_2) are *conjugate* if there is a homeomorphism $h : X_1 \rightarrow X_2$ such that $h \circ \sigma_1 = \sigma_2 \circ h$.

Shifts of finite type

Definition 1

A shift space X over a finite alphabet \mathcal{A} is of *finite type* if it may be described using a finite list of forbidden words.

Example 2

Let $E = (E^0, E^1, r, s)$ be a row-finite directed graph with no sinks or sources and E^0 finite. Let

$$X_E = \{x \in (E^1)^{\mathbb{Z}} : r(x_i) = s(x_{i+1})\}$$

then X_E is a shift of finite type with $\mathcal{F} = \{ef : r(e) \neq s(f)\}$. The shift space X_E is called the *edge shift* associated to E .

We differ here from the usual convention in symbolic dynamics that X_E consists of bi-infinite paths in E running from left to right.

Theorem 3 (Folklore)

Let X be a shift of finite type, then there is a directed graph E such that X is conjugate to X_E .

Conjugacy of shifts

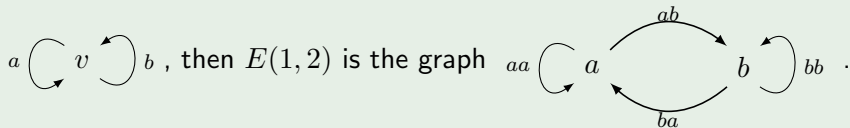
One of the most basic examples of a conjugacy is to take shift space and create a new shift using an alphabet which consists of words of a fixed length in the original shift, this process is called a *higher-block code*, which can easily be seen for shifts of finite type.

Example 4

Let $E = (E^0, E^1, r, s)$ for $m \geq 0$ let $E(m, m+1) = (E^m, E^{m+1}, r, s)$ where

$$r(\alpha_1 \dots \alpha_{m+1}) = \alpha_2 \dots \alpha_{m+1} \text{ and } s(\alpha_1 \dots \alpha_{m+1}) = \alpha_1 \dots \alpha_m.$$

Then X_E is conjugate to $X_{E(m, m+1)}$. For instance if E is the graph



Congruency and graph C^* -algebras

Theorem 5 (see [1])

Let E be a row-finite graph with no sources and $m \geq 0$. Let (s, p) be a Cuntz-Krieger E -family generating $C^*(E)$. For $\alpha \in E^m$ and $\beta \in E^{m+1}$ set $T_\beta := s_\beta s_{\beta_2 \dots \beta_{m+1}}^*$, $Q_\alpha = s_\alpha s_\alpha^*$. Then $\{T, Q\}$ is a Cuntz-Krieger $E(m, m+1)$ family in $C^*(E)$ and the map $\pi_{T, Q} : C^*(E(m, m+1)) \rightarrow C^*(E)$ is an isomorphism.

Example 6

If E is the graph $a \begin{pmatrix} \curvearrowright v \curvearrowright \end{pmatrix} b$, then $C^*(E) \cong C^*(E(1, 2))$ where $E(1, 2)$ is

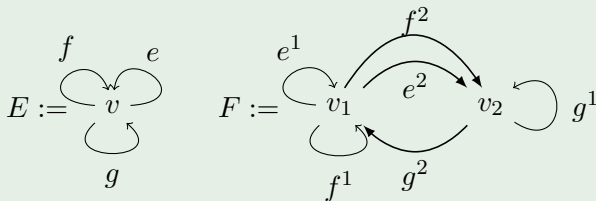
the graph $aa \begin{pmatrix} \curvearrowright a \end{pmatrix} \begin{matrix} \xrightarrow{ab} \\ \xleftarrow{ba} \end{matrix} b \begin{pmatrix} \curvearrowright bb \end{pmatrix}$.

Conjugacy and graph C^* -algebras II

- It can be shown that the shifts of finite type X_E and X_F associated to graphs E and F are conjugate if and only if there is a finite sequence of certain graphical moves transforming E into F (see [10]).
- In [2] (see also [8]) it was shown how these moves affect the corresponding graph C^* -algebras, and for more general moves in [12].

Example 7

The graphs E, F shown below are related via an out-splitting procedure, where vertex v splits into v_1 , which emits e, f and v_2 which emits g ; and so $C^*(E) \cong C^*(F)$.



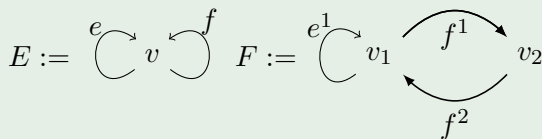
Flow equivalence of shift spaces

There is a weaker notion of equivalence between shift spaces called *flow-equivalence*. For shift of finite type it reduces to inserting (or deleting) a vertex in an edge, a graphical move called a *delay*.



Example 8

The graph F is obtained from E by inserting a vertex in the edge f . So X_E and X_F are flow equivalent.

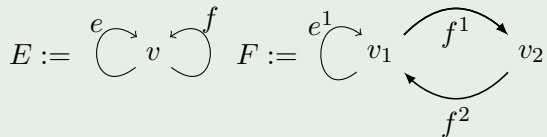


Flow equivalence and graph C^* -algebras

In [2] and [6] (see also [12]) it was shown that if F is obtained from E by a delay then $C^*(E)$ and $C^*(F)$ are Morita equivalent.

Example 9

For the graphs E, F below, $C^*(E)$ and $C^*(F)$ are Morita equivalent since the graph F is obtained from E by a delay.



Invariants for shift spaces

Let E be a row-finite graph with no sources and r vertices. Define the $r \times r$ *vertex matrix* A_E by

$$A_E(v, w) = \#\{e \in E^1 : s(e) = v, r(e) = w\}.$$

Then the *Bowen-Franks* group associated to the shift of finite type (X_E, σ_E) is

$$BF(X_E) = \text{coker}(1 - A_E^t : \mathbb{Z}^r \rightarrow \mathbb{Z}^r) := \mathbb{Z}^r / \text{Im}(1 - A_E^t)\mathbb{Z}^r.$$

In [5] Bowen and Franks showed that this abelian group is an invariant for conjugacy and flow equivalence of shifts of finite type.

Invariants for graph C^* -algebras

For a C^* -algebra A there are two abelian groups $K_0(A), K_1(A)$ which are invariants for Morita equivalence. Roughly speaking, $K_0(A)$ measures the dimension of projections in A and $K_1(A)$ measures the connectivity of the unitaries in A .

Theorem 10 (see [7], [4])

Let E be a row-finite directed graph with no sinks or sources and r vertices. Then

$$\begin{aligned} K_0(C^*(E)) &\cong \operatorname{coker}(1 - A_E^t : \mathbb{Z}^r \rightarrow \mathbb{Z}^r), \\ K_1(C^*(E)) &\cong \ker(1 - A_E^t : \mathbb{Z}^r \rightarrow \mathbb{Z}^r). \end{aligned}$$

In particular we note that $K_0(C^*(E))$ is the Bowen-Franks group for the shift of finite type X_E .

Higher rank graphs

Definition 11 (see [9])

Let $k \geq 1$. A higher rank graph, or *k -graph*, is a countable category Λ equipped with a degree functor $d : \Lambda \rightarrow \mathbb{N}^k$ satisfying the *factorisation property*: if $d(\lambda) = m + n$ then there exist unique μ and ν such that $d(\mu) = m$, $d(\nu) = n$ and $\lambda = \mu\nu$.

- For $k = 1$, the category Λ is the path category of a directed graph, and d is the length function.
- For $k > 1$ we think of Λ as a certain collection of paths in a multicoloured directed graph.

Higher rank graphs II

- In [9] it is shown how associate a C^* -algebra to a k -graph.
- The class of k -graph C^* -algebras is larger than the class of 1-graph algebras, containing the $A\mathbb{T}$ algebras (see [11]) for $k \geq 2$.
- Many results which hold for directed graphs (1-graphs) but the arguments are often more subtle.
- Unfortunately some results for k -graph C^* -algebras are not so clear cut: for instance there is no useful formula for K -theory for $k > 2$.

Labelled graphs

A *labelled graph* (E, \mathcal{L}) over a countable alphabet \mathcal{A} consists of a directed graph E together with a labelling map $\mathcal{L} : E^1 \rightarrow \mathcal{A}$. Without loss of generality we may assume that the map \mathcal{L} is onto.

For $x \in X_E$ define $\mathcal{L}(x) \in \mathcal{A}^{\mathbb{Z}}$ by

$$\mathcal{L}(x)_i = \mathcal{L}(x_i) \text{ for } i \in \mathbb{Z}$$

A labelled graph (E, \mathcal{L}) gives rise to a shift space $(X_{(E, \mathcal{L})}, \sigma)$ by

$$X_{(E, \mathcal{L})} = \{y \in \mathcal{A}^{\mathbb{Z}} : y = \mathcal{L}(x) \text{ for some } x \in X_E\}.$$

The labelled graph (E, \mathcal{L}) is called a *presentation* of the shift space X if $X = X_{(E, \mathcal{L})}$.

If \mathcal{L} is injective, then (E, \mathcal{L}) may be regarded as a directed graph. Every shift space may be presented by a labelled graph.

Labelled graphs II

- The C^* -algebra $C^*(E, \mathcal{L})$ associated to a labelled graph (E, \mathcal{L}) was introduced in [3].
- If the labelling map $\mathcal{L} : E^1 \rightarrow \mathcal{A}$ is injective then $C^*(E, \mathcal{L})$ is isomorphic to $C^*(E)$.
- If E^0 is finite then the shift space $X_{(E, \mathcal{L})}$ is called *sofic*. In this case it was shown that in [3] that $C^*(E, \mathcal{L}) \cong C^*(E)$.
- There are examples of labelled graph C^* -algebras which are unital but have infinitely many vertices and countably generated K_0 -groups.
- The relationship between graph C^* -algebras and labelled graph C^* -algebras is not fully understood.

References I

- ① T. Bates, *Applications of the gauge-invariant uniqueness theorem*, Bull. Austral. Math. Soc., **66** (2002), 57–67.
- ② T. Bates and D. Pask, *Flow equivalence of graph algebras*, Ergodic Th. & Dynam. Sys. **24** (2004), 367–382.
- ③ T. Bates and D. Pask, *C^* -algebras of labelled graphs*, J. Operator Theory **57** (2007), 207–226.
- ④ T. Bates, D. Pask, I. Raeburn and W. Szymański, *The C^* -algebras of row-finite graphs*, New York J. Math. **6** (2000), 307–324.
- ⑤ R. Bowen and J. Franks, *Homology for zero-dimensional nonwandering sets*, Ann. Math. **106** (1977), 73–92.
- ⑥ T. Crisp and D. Gow, *Contractible subgraphs and Morita equivalence of graph C^* -algebras*, Proc. Amer. Math. Soc. **135** (2006), 2003–2013.
- ⑦ J. Cuntz, *A class of C^* -algebras and topological Markov chains II: Reducible chains and the Ext-functor for C^* -algebras*, Invent. Math., **63** (1981) 25–40.
- ⑧ D. Drinen and N. Sieben, *C^* -equivalences of graphs*, J. Operator Theory, **45** (2001), 209–229.

References II

- 9 A. Kumjian and D. Pask, *Higher rank graph C^* -algebras*, New York J. Math. **6** (2000), 1–20.
- 10 D. Lind and B. Marcus, *An introduction to symbolic dynamics and coding*. Cambridge University Press., 1995.
- 11 D. Pask, I. Raeburn, M. Rørdam and A. Sims, *Rank-Two Graphs whose C^* -algebras are direct limits of Circle Algebras*, *J. Funct. Anal.*, **239** (2006), 137–178.
- 12 A. Sørensen, *Geometric classification of simple graph algebras*, arXiv:1111.1592).